

# Buffer Overflow Attacks

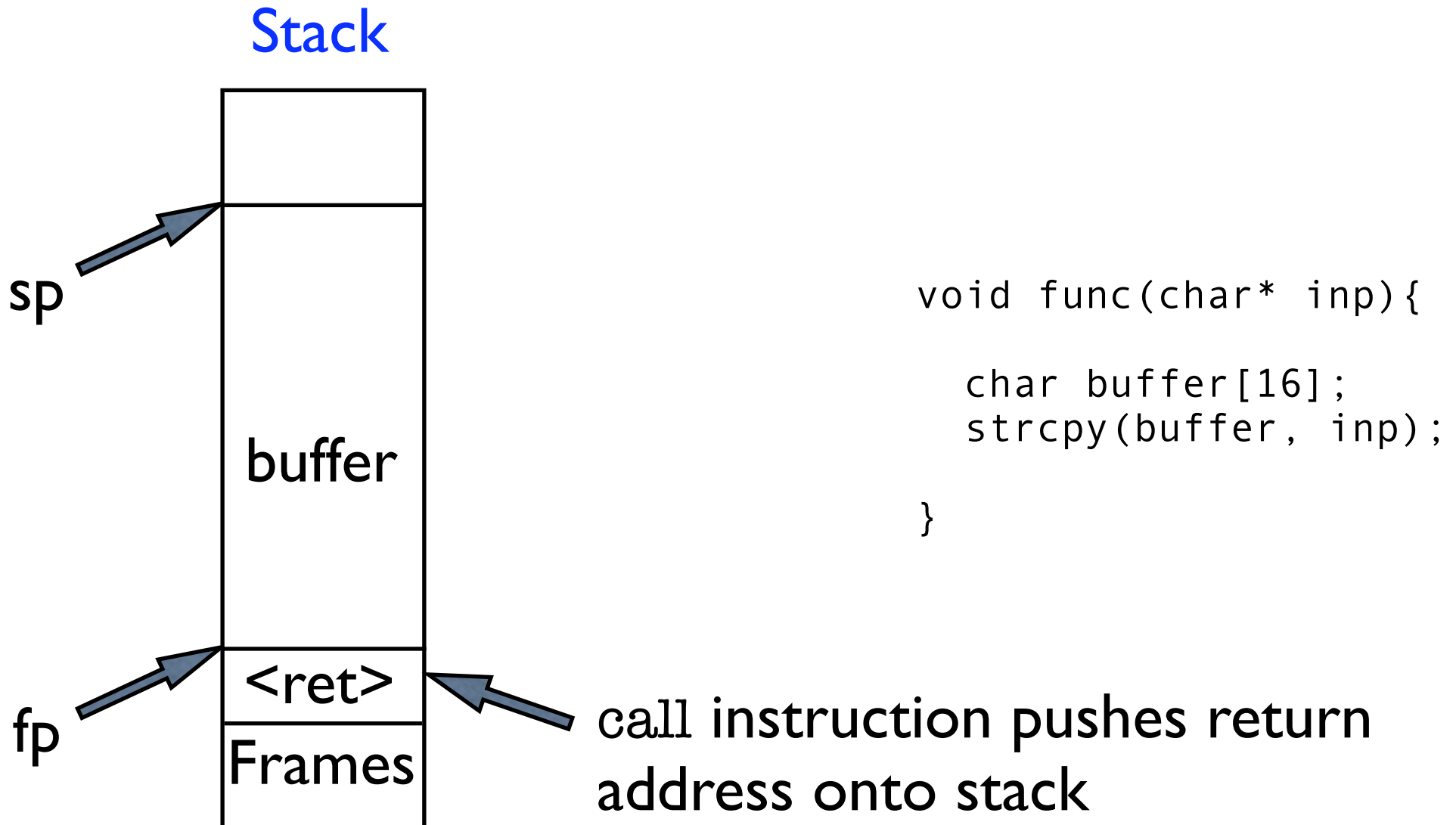
1. Smashing the Stack
2. Other Buffer Overflow Attacks
3. Work on Preventing Buffer Overflow Attacks

# Smashing the Stack

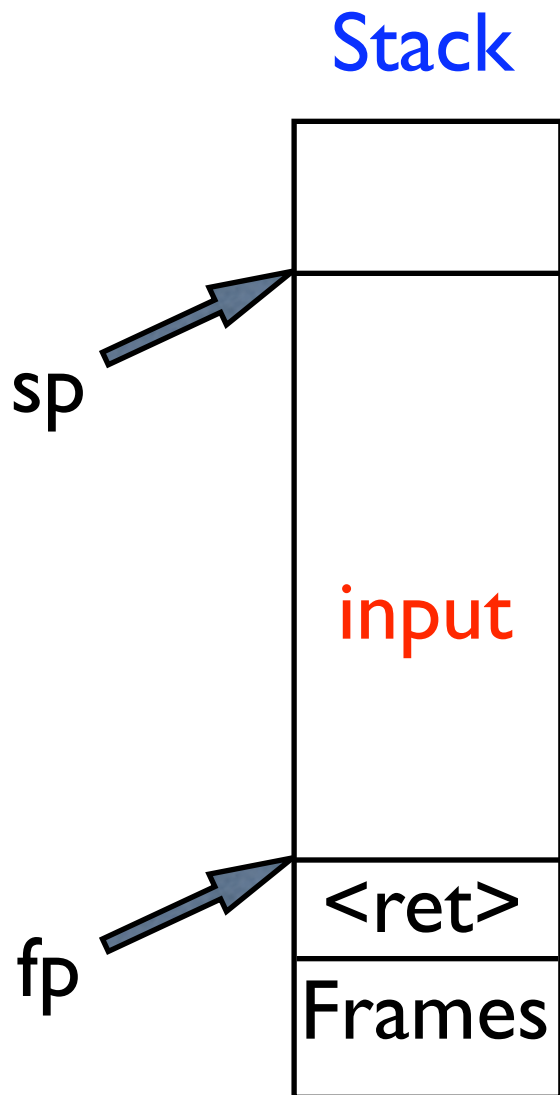
# An Evil Function

```
void func(char* inp) {  
    char buffer[16];  
    strcpy(buffer, inp);  
}
```

# Calling func

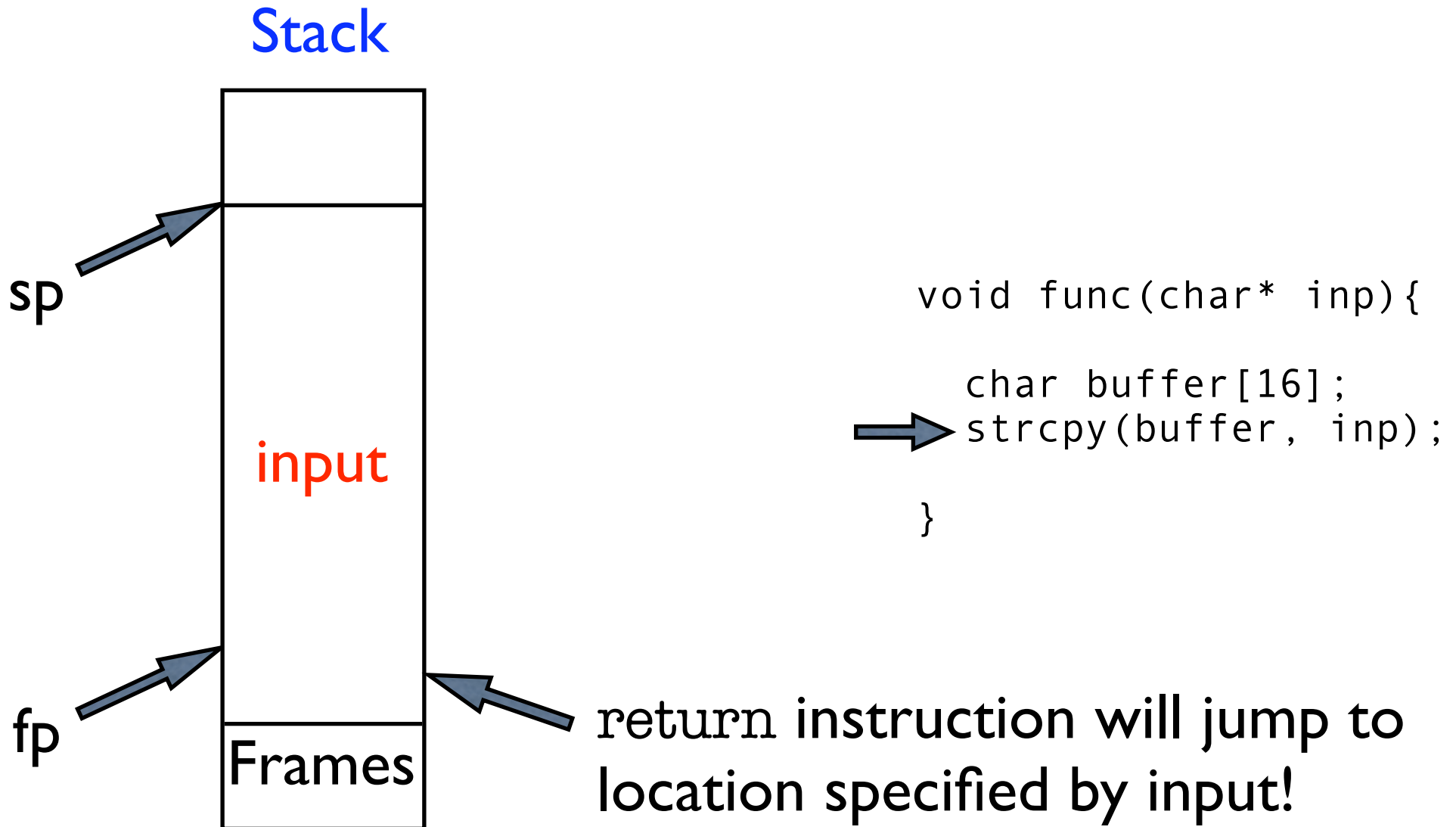


# Operation of func

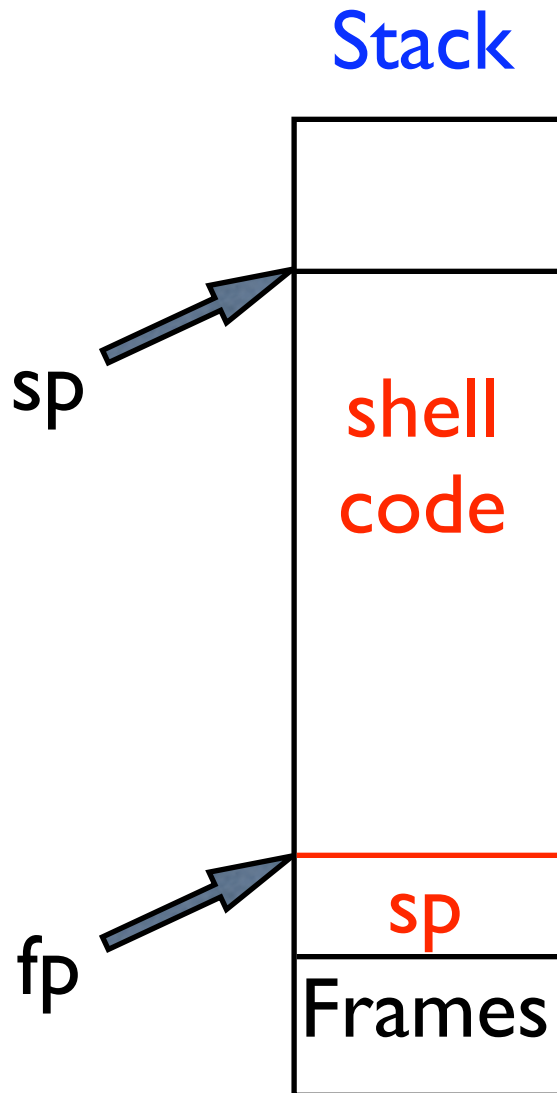


```
void func(char* inp){  
    char buffer[16];  
    → strcpy(buffer, inp);  
}
```

# C doesn't check bounds...



# The Exploit



```
void func(char* inp){  
    char buffer[16];  
    strcpy(buffer, inp);  
}
```



# Two Questions

1. How does attacker know where buffer starts on the stack?

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2. How does attacker know where to put ra in the input to overwrite <ret>?

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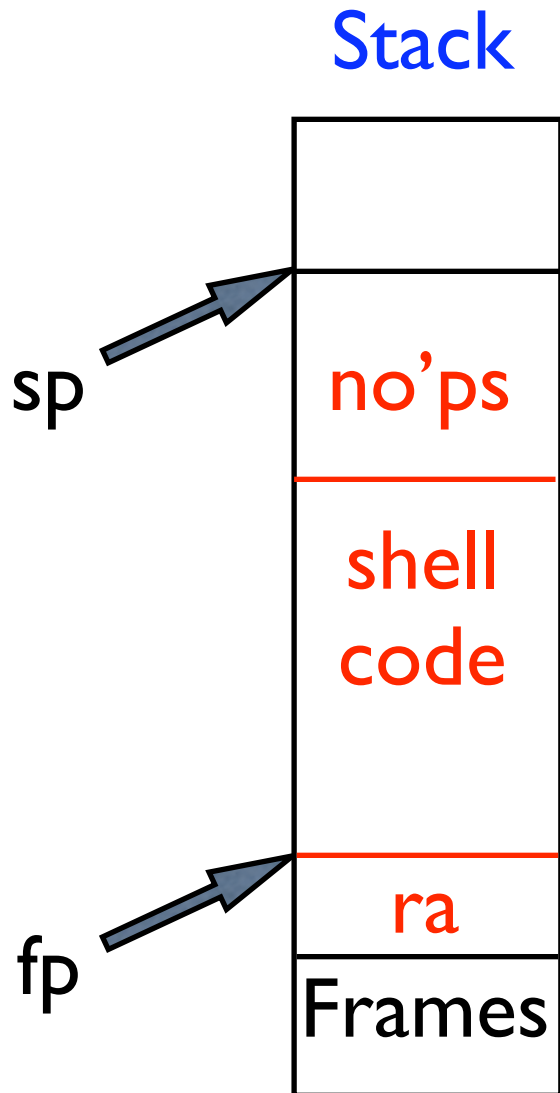
1. How does attacker know where buffer starts on the stack?
2. How does attacker know where to put ra in the input to overwrite <ret>?

Answer: she doesn't, but she can guess...

# Guessing ra

- Stack has same starting location in all programs
- Most programs don't ever put more than a few thousand bytes on the stack
- So, just guess ra!

# Adding a “Launching Pad”

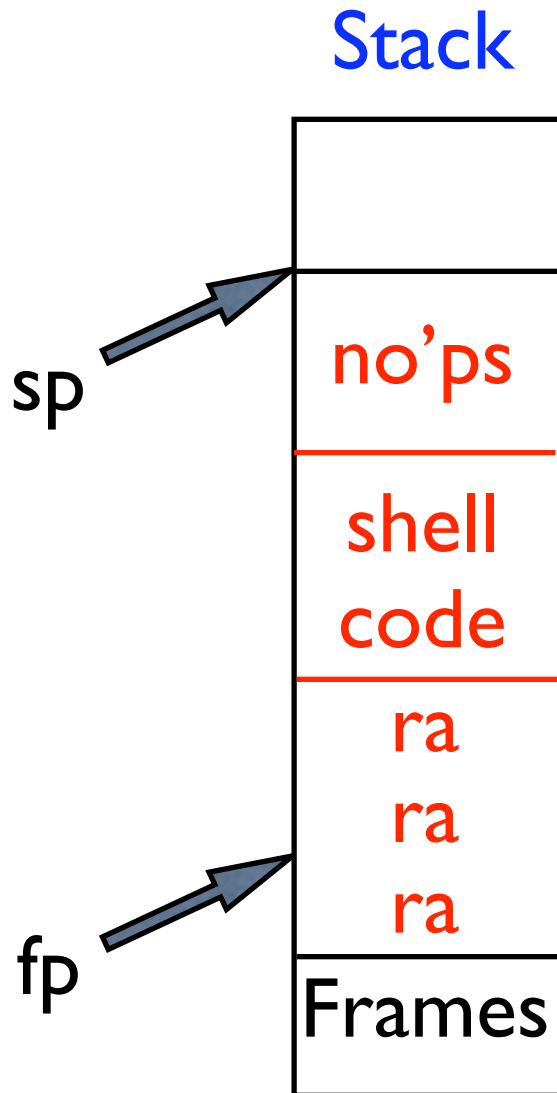


```
void func(char* inp){  
    char buffer[16];  
    strcpy(buffer, inp);  
}
```

# Guessing Where to Put ra in the Input

- This is really just trial-and-error...
- But we can improve our chances on this as well.

# Improvement



```
void func(char* inp){  
    char buffer[16];  
    strcpy(buffer, inp);  
}
```

# Final Version of Input

no'ps	shellcode	ra ra ra
-------	-----------	----------

A good value of ra is worked out through trial and error, as is a good length for the input



# Other Buffer Overflow Attacks

# Arc Injection

- Input overwrites return address to point to a piece of code *not* supplied by the attacker
- For example, system call in libc:

```
void system(char * arg){  
    check_validity(arg);  
    R = arg;  
    execl(R, ...) }  
register          target for attacker
```

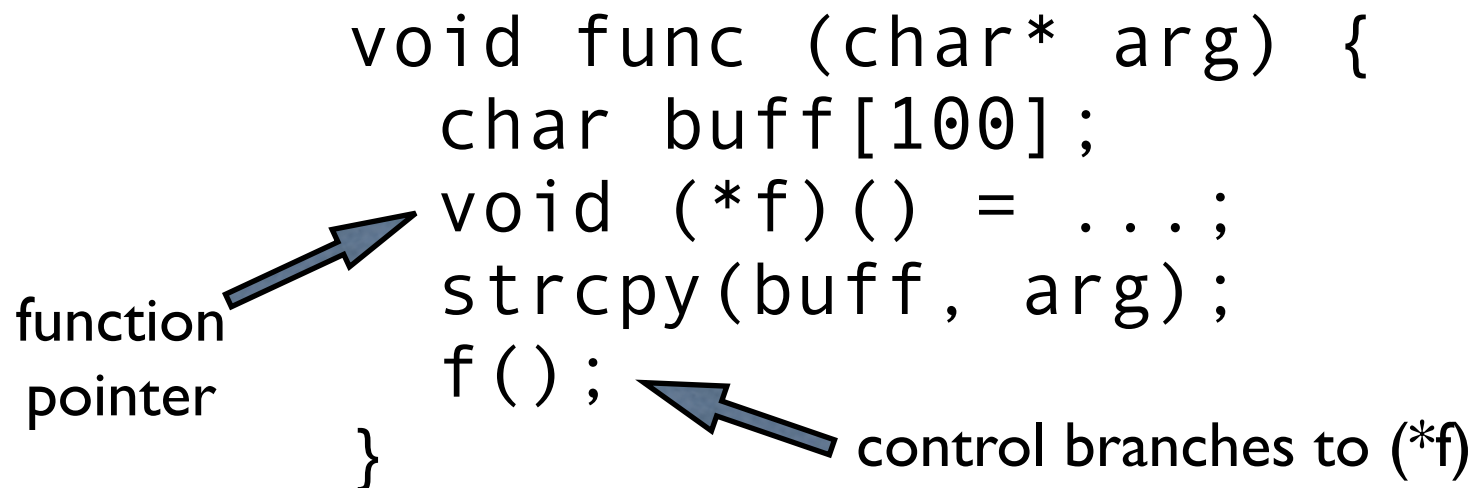
# Pointer Subterfuge

- Attacker overwrites a pointer instead of the return address
- Example: overwriting a function pointer

```
void func (char* arg) {  
    char buff[100];  
    void (*f)() = ...;  
    strcpy(buff, arg);  
    f();  
}
```

function  
pointer

control branches to (\*f)

The diagram illustrates a function pointer subterfuge. It shows a C function named 'func' that takes a character pointer 'arg' and returns void. Inside the function, a character array 'buff' of size 100 is declared. A function pointer 'f' is declared as 'void (\*f)()' and assigned a value represented by '...'. The function then calls 'strcpy(buff, arg);' and finally calls 'f();'. An arrow labeled 'function pointer' points to the declaration of 'f'. Another arrow labeled 'control branches to (\*f)' points to the call 'f();', indicating that the program's execution flow is controlled by the function pointer 'f'.

# Heap Smashing

- Exploits buffer overflows in heap rather than on stack
- Often combined with pointer subterfuge (eg, overwriting a function pointer that is not a local variable)
- Gaining in popularity

# Preventing Buffer Overflows

# Stackguard (Cowan et al.)

- Instruments code during compilation:
  1. Augments call instructions to write a 'canary word' next to return address, and push it onto a 'canary stack'
  2. Augments return instructions to first check that canary word is intact

# ProPolice (Etoh & Yoda)

- Extension to Stackguard's method: they place canary word between local `char[]`'s and other local variables
- This protects against pointer subterfuge on the stack

# StackShield

- Similar to StackGuard, but saves a copy of return address to check against
- If real return address has been overwritten, can either terminate program or use copy and let program keep running



# CRED: Bounds Checking for C (Ruwase & Lam)

- Does bounds checking for all string/char[] accesses
- Only tool to detect all attacks in a testbed that includes all previously mentioned attacks
- Adds 26%-130% overhead to the program

# Baratloo et al.

- Two dynamically loaded libraries: Libsafe replaces unsafe library functions with safe versions, and Libverify does return address verification
- Works on executables
- Slowdown is  $< 15\%$  for all tested programs

# Baratloo et al.

- Weaknesses of Libverify:
  1. Only prevents attacks that overwrite the return address
  2. For technical reasons, has to copy entire function to heap on each call

# References

- Aleph I. “Smashing the stack for fun and profit.”  
<http://www.phrack.org/phrack/49/P49-16>
- Baratloo, Singh, and Tsai. “Transparent runtime defense against smash stacking attacks.”  
[http://www.usenix.org/publications/library/proceedings/usenix2000/general/full\\_papers/baratloo/baratloo.pdf](http://www.usenix.org/publications/library/proceedings/usenix2000/general/full_papers/baratloo/baratloo.pdf)
- Cowan et al. “StackGuard”  
<http://www.cse.ogi.edu/DISC/projects/immunix/StackGuard/usenixsc98.ps.gz>

# References

- Etoh & Yoda. “Protecting from stack-smashing attacks.”  
<http://www.trl.ibm.com/projects/security/ssp/main.html>
- Ruwase & Lam. “A practical dynamic buffer overflow detector.”  
<http://suif.stanford.edu/papers/tunji04.pdf>
- Stackshield.  
<http://www.angelfire.com/sk/stackshield>